

# ECE 2400 Computer Systems Programming

## Spring 2025

### Topic 6: C Dynamic Allocation

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Cornell University

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1	Using <code>malloc</code> to Allocate Memory	2
2	Using <code>free</code> to Deallocate Memory	8
3	Mapping Conceptual Storage to Machine Memory	10

**zyBooks** logo indicates readings and coding labs in the course zyBook which will not be discussed in detail in lecture. Students are responsible for all material covered in lecture and in the course zyBook.

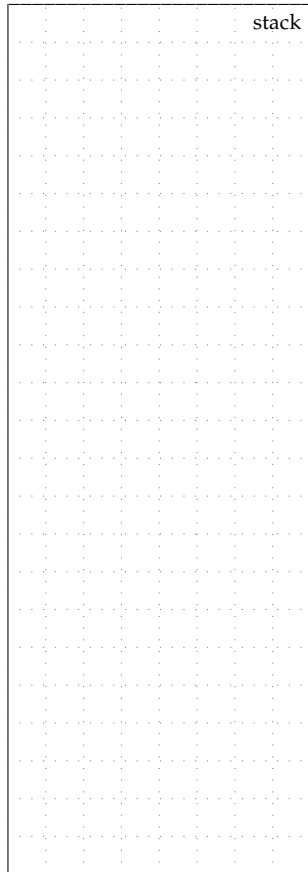
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## 1. Using malloc to Allocate Memory

- Let's revisit an example we saw in a previous topic
- Assume we wish to refactor prepending a node to the front of a chain of nodes into its own function

Draw a state diagram corresponding to the execution of this program

```
0001 typedef struct _node_t
0002 {
0003     int          value;
0004     struct _node_t* next_p;
0005 }
0006 node_t;
0007
0008 node_t* prepend( node_t* head_p,
0009                 int v )
0010 {
0011     node_t node;
0012     node.value = v;
0013     node.next_p = head_p;
0014     return &node;
0015 }
0016
0017 int main( void )
0018 {
0019     node_t* head_p = NULL;
0020     head_p = prepend( head_p, 3 );
0021     head_p = prepend( head_p, 4 );
0022     return 0;
0023 }
```



- Let's consider a similar idea for arrays
- Assume we wish to refactor allocating an array and then initializing all elements to zero into its own function

```
1  int* init_array( int n )
2  {
3      int x[n];
4
5      for ( int i=0; i<n; i++ )
6          x[i] = 0;
7
8      return x;
9  }
10
11 int main( void )
12 {
13     int* a = init_array(3);
14     return 0;
15 }
```

**List two errors with this function:**

1. \_\_\_\_\_

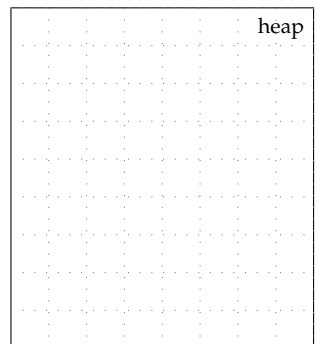
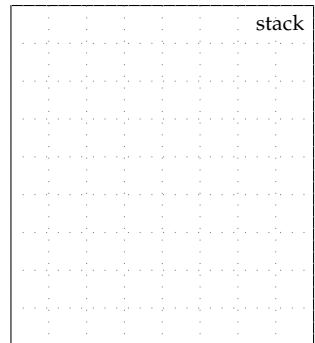
\_\_\_\_\_

2. \_\_\_\_\_

\_\_\_\_\_

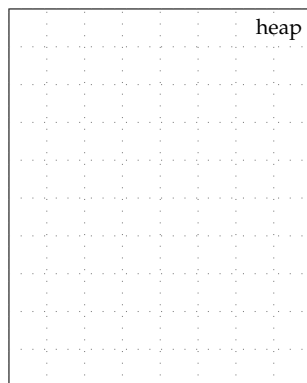
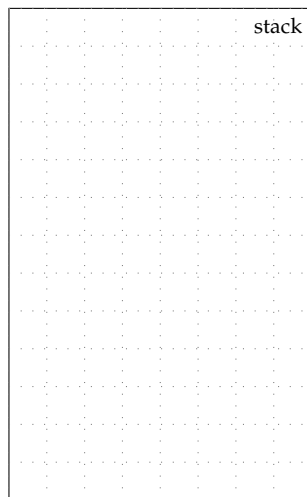
- **Dynamic memory allocation** uses the **heap** (new region of memory)
- Because dynamically allocated variables are not on a function's stack frame, they are not deallocated when a function returns
- We can dynamically allocate variables on the heap using `malloc`
- `malloc` takes the number of bytes to allocate as a parameter and returns a pointer to the new variable allocated on the heap
- Since the amount of memory allocated is dynamic, we can create arrays where the number of elements is not known until runtime
- `malloc` is defined in `stdlib.h`

```
□□□ 01 int* a_ptr =
□□□ 02     malloc( sizeof(int) );
□□□ 03
□□□ 04 *a_ptr = 42;
□□□ 05
□□□ 06 int* b_ptr =
□□□ 07     malloc( 4 * sizeof(int) );
□□□ 08
□□□ 09 b_ptr[0] = 10;
□□□ 10 b_ptr[1] = 11;
□□□ 11 b_ptr[2] = 12;
□□□ 12 b_ptr[3] = 13;
```



Draw a state diagram corresponding to the execution of this program

```
□□□ 01 typedef struct
□□□ 02 {
□□□ 03     double real;
□□□ 04     double imag;
□□□ 05 }
□□□ 06 complex_t;
□□□ 07
□□□ 08 int main( void )
□□□ 09 {
□□□ 10     complex_t* c_ptr0 =
□□□ 11         malloc( sizeof(complex_t) );
□□□ 12
□□□ 13     c_ptr0->real = 1.5;
□□□ 14     c_ptr0->imag = 3.5;
□□□ 15
□□□ 16     complex_t* c_ptr1 =
□□□ 17         malloc( sizeof(complex_t) );
□□□ 18
□□□ 19     c_ptr1->real = c_ptr0->real;
□□□ 20     c_ptr1->imag = c_ptr0->imag;
□□□ 21
□□□ 22     return 0;
□□□ 23 }
```

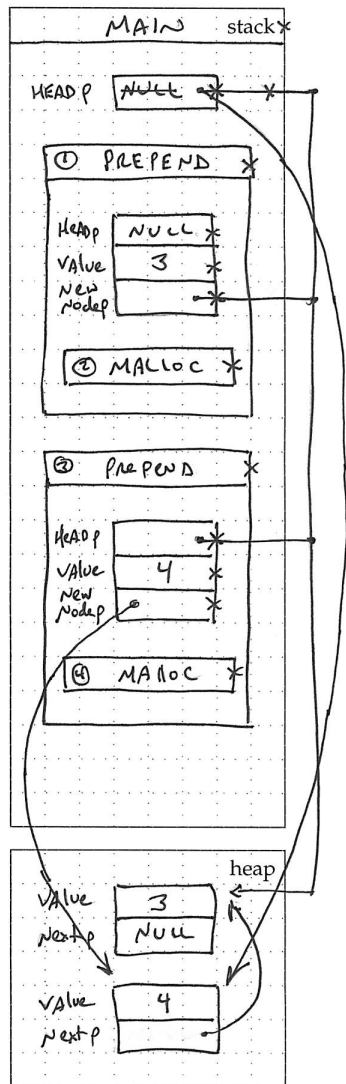


## 1. Using malloc to Allocate Memory

- Assume we wish to refactor prepending a node to the front of a chain of nodes into its own function

```
01 typedef struct _node_t
02 {
03     int         value;
04     struct _node_t* next_p;
05 }
06 node_t;
07
08 node_t* prepend( node_t* head_p,
09                 int v )
10 {
11     node_t* new_node_p =
12         malloc( sizeof(node_t) );
13
14     new_node_p->value = v;
15     new_node_p->next_p = head_p;
16     return new_node_p;
17 }
18
19 int main( void )
20 {
21     node_t* head_p = NULL;
22     head_p = prepend( head_p, 3 );
23     head_p = prepend( head_p, 4 );
24     return 0;
25 }
```

**zyBooks** The course zyBook includes a coding lab to implement a function to append a node to the *back* of a chain of nodes.



- Assume we wish to refactor allocating an array and then initializing all elements to zero into its own function

```
1  int* init_array( int n )
2  {
3      int* x = malloc( n * sizeof(int) );
4
5      for ( int i=0; i<n; i++ )
6          x[i] = 0;
7
8      return x;
9  }
10
11 int main( void )
12 {
13     int* a = init_array(3);
14     return 0;
15 }
```

**How does this address the two errors we identified earlier?**

1. \_\_\_\_\_

\_\_\_\_\_

2. \_\_\_\_\_

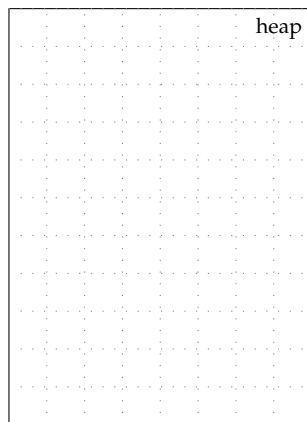
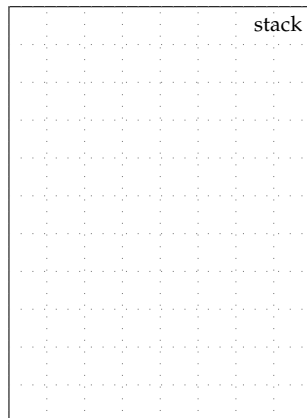
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**zyBooks** The course zyBook includes a coding lab to implement a function to duplicate a given array on the heap and return a pointer to this newly allocated array.

## 2. Using free to Deallocate Memory

Draw a state diagram corresponding to the execution of this program

```
□□□ 01 typedef struct
□□□ 02 {
□□□ 03     double real;
□□□ 04     double imag;
□□□ 05 }
□□□ 06 complex_t;
□□□ 07
□□□ 08 int main( void )
□□□ 09 {
□□□ 10     complex_t* c_ptr =
□□□ 11         malloc( sizeof(complex_t) );
□□□ 12
□□□ 13     c_ptr->real = 1.5;
□□□ 14     c_ptr->imag = 3.5;
□□□ 15
□□□ 16     c_ptr =
□□□ 17         malloc( sizeof(complex_t) );
□□□ 18
□□□ 19     c_ptr->real = 2.5;
□□□ 20     c_ptr->imag = 4.5;
□□□ 21
□□□ 22     return 0;
□□□ 23 }
```

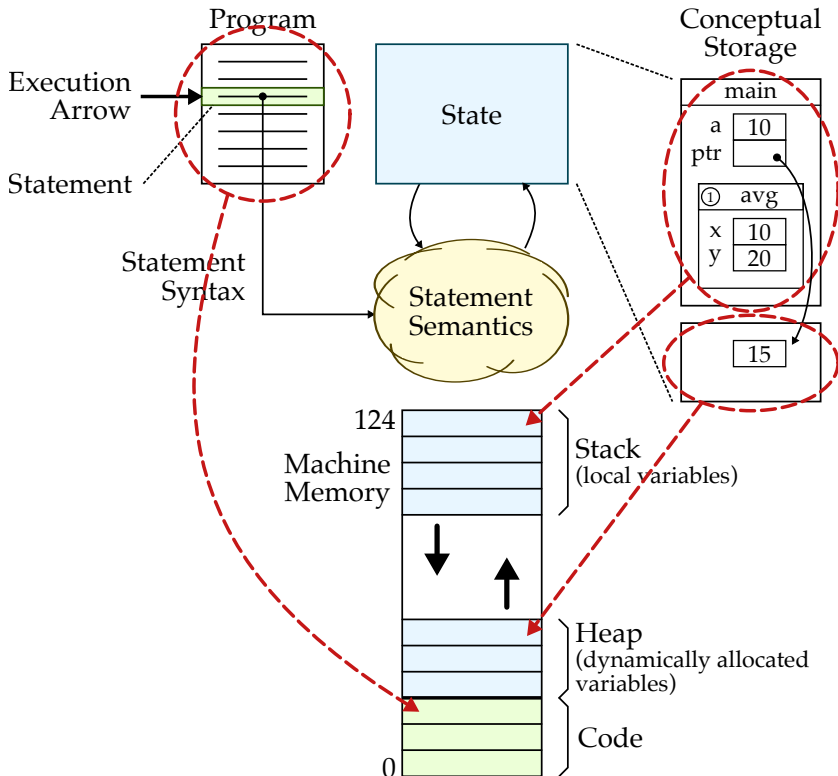


- Every call to malloc must have corresponding call to free
- free takes a pointer to a dynamically allocated variable

```
1  typedef struct
2  {
3      double real;
4      double imag;
5  }
6  complex_t;
7
8  int main( void )
9  {
10     complex_t* c_ptr =
11         malloc( sizeof(complex_t) );
12
13     c_ptr->real = 1.5;
14     c_ptr->imag = 3.5;
15
16     free( c_ptr );
17
18     c_ptr =
19         malloc( sizeof(complex_t) );
20
21     c_ptr->real = 2.5;
22     c_ptr->imag = 4.5;
23
24     free( c_ptr );
25
26     return 0;
27 }
```

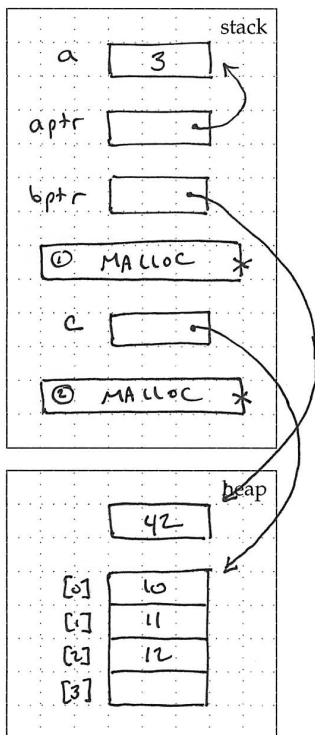
### 3. Mapping Conceptual Storage to Machine Memory

- Recall that our current use of state diagrams is conceptual
- Real machine uses **memory** to store variables
- Real machine does not use “arrows”, uses **memory addresses**
- Heap is stored above code and grows *up*

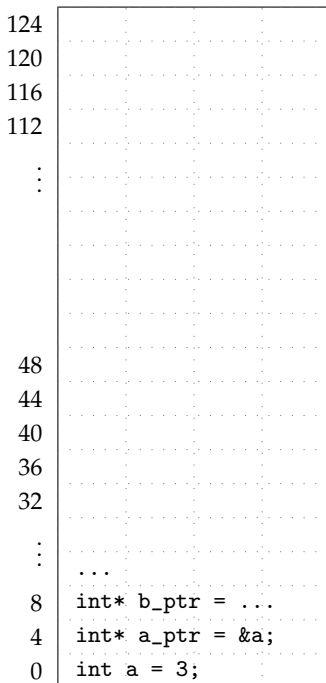


### 3. Mapping Conceptual Storage to Machine Memory

```
□□□ 01 int a = 3;
□□□ 02 int* a_ptr = &a;
□□□ 03
□□□ 04 int* b_ptr = malloc( sizeof(int) );
□□□ 05 *b_ptr = 42;
□□□ 06
□□□ 07 int* c = malloc( 4 * sizeof(int) );
□□□ 08 c[0] = 10;
□□□ 09 c[1] = 11;
□□□ 10 c[2] = 12;
```



**Memory**  
(4B word addr)



## Machine memory in real systems

- Machine memory size ranges from KBs (embedded) to TBs (server)
- Lowest address range reserved to detect NULL pointer dereference
- Static data region is used for global variables
- Machine memory as shown is really the *virtual memory space*
- Different programs have their own virtual memory spaces mapped to a single large *physical memory space*

